

A Genetic Algorithm for Solving Periodic Heterogeneous Vehicle Routing Problem

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ABSTRACT

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Keywords Periodic Routing Problem Heterogeneous Vehicles Genetic Algorithm Taguchi This paper addresses the periodic heterogeneous vehicle routing problem (PHVRP), an extension of the classical vehicle routing problems (VRP). This problem is known to be confined to various real-world instances where each customer's demand should be served within a specific time horizon and a maximum demand quantity that can be delivered at each visit. The heterogeneous capacitated vehicles are available to perform the services for each customer. This paper aims to minimize the total traveling time of routes for all vehicles over the time horizon so that the customers' demands can be delivered. Thus, a novel coding scheme is also proposed to directly convert a random sequence of integers into a feasible solution, which is then embedded into algorithms. Furthermore, this paper also compares the performance of the Genetic Algorithm (GA) with the particle swarm optimization algorithm (PSO). The numerical results of the experiments show that the proposed GA is superior to PSO. However, the computation time of PSO is faster than GA.

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1. Introduction

A basic vehicle routing problem (VRP) consists of a set of customers, each with a fixed quantity of demand, which must be served by a vehicle that travels from a depot to the customer and then returns to the depot [1-3]. Delivery routes of a vehicle must begin and end at the depot to ensure that all customer demands are satisfied and that each customer is visited by only one vehicle. Vehicle capacities are specified, and each vehicle frequently has a maximum total traveling length. Toth and Vigo [4] deal with the vehicle routing problem by creating a distribution scheduling plan in which vehicles are assigned to pick up and deliver products. The problem involves a single period, meaning that customer visit time starts at the beginning of the visit period and terminates when all customers have been served. Additionally, some real-life implementations are concerned with the periodic delivery operations during a specific time horizon, i.e., liquid petroleum of gasoline delivery, grocery delivery, or waste collections. According to Campbell and Wilson [5], periodic vehicle routing problems (PVRP) are an extension of the VRP in which



customers should be served at a particular frequency of visit according to a set schedule and receive a defined demand at each customer visit. The PVRP expands a single delivery period to T periods (T > 1), allowing each customer to require several services over the T period, i.e., within a one-week or one-month delivery period. Generally, VRP research focuses on reducing travel distances and times [6, 7]. Similarly, the PVRP needs to optimize each delivery period for a specific customer combination and the path selection for each delivery period to reduce the total traveling distance. Therefore, the PVRP is more complex than a classical VRP.

Beltrami and Bodin first applied PVRP to assign compactor trucks to a city waste collection [8]. Afterward, Chao, et al. [9] developed a heuristic method to solve the PVRP using the initialization and improvement phases. Baptista, et al. [10] presented an improved heuristic for solving the PVRP in waste paper box recycling. Francis, et al. [11] experimented with an improved heuristic algorithm for solving the PVRP using customer service frequencies as decision variables. Pourghaderi, et al. [12] developed a new simple, effective line construction heuristic algorithm for solving the PVRP using an embedded improvement procedure. Matos and Oliveira [13] experimented with the ant colony algorithm by suggesting a novel strategy to update the pheromone. Yu and Yang [14], proposed an ant colony algorithm based on two crossover operations to solve the PVRPTW. Garside and Laili [15] solved the periodic multi-trip vehicle routing problem using a cluster first route second heuristic. They found that the cluster first route second could generate a better routing solution. Some other heuristics for PVRP can be found in subsequent research [16-18].

This paper studies the periodic heterogeneous vehicle routing problem (PHVRP) motivated by the real-life scheduling of vehicles involving different vehicle capacities and periodic deliveries. In order to complete PHVRP, each delivery period must be optimized, as well as each delivery period using various vehicle capacities. Therefore, the PHVRP is embedded in the classical VRP. It is more complex since it includes an NP-Hard problem and can be solved using the exact and approximate method. According to researchers, there needs to be more literature relevant to the PHVRP. One related study is to develop an efficient technique for solving the periodic heterogeneous vehicle routing problem with driver scheduling using a heuristic algorithm and exact method [19]. Vidal, et al. [20] solved the multi-depot periodic vehicle routing problem in which the company has more than one depot and each depot has a different vehicle using a genetic algorithm. Furthermore, Cantu-Funes, et al. [21] solved the problem of the multi-depot periodic vehicle routing concerning the due dates and time windows. They solved the problem by using a mixed integer method and a reactive greedy randomized adaptive search procedure and determining the heterogeneous vehicles, the schedule of each vehicle, and the route between the factory and warehouse. Yao, et al. [22] solved the PVRP with different capacities of vehicles for each test problem. They developed an artificial bee colony algorithm by using a scanning strategy.

One of the approximate methods is the genetic algorithm (GA). GA is an evolutionary algorithm (EA) inspired by the natural selection mechanism introduced and initiated in 1970 by Holland. Researchers have drawn great attention to the GA due to its robustness and flexibility. Thus, the genetic algorithm has been used to solve various combinatorial problems, including certain types of vehicle routing difficulties [23]. This algorithm generally searches a very large and potentially high-dimensional search space. The genetic algorithm is a powerful component in a random search to solve uncertain and complex problems requiring a large time-space to find the optimal solution [24]. It is an effective process for finding a great solution set available for a specific design. Several researchers have successfully used the genetic algorithm in their studies [25-28]. The

genetic algorithm was developed to solve the multi-depot vehicle routing problem Mirabi [29], where the hybrid genetic algorithm competes with previous methods regarding problem-solution quality. Nguyen, et al. [30] proposed the GA with two neighborhoods to solve the PVRP with time windows.

The goal of this paper is to minimize the total length traveled of routes for all vehicles during the time horizon so that customer demands are served. The rest of this paper is organized as follows. Section 2 describes the characteristic system, followed by a novel coding procedure, a description of the PVRP, the genetic algorithm, and the experimental setup using the Taguchi method. Section 3 presents the numerical and computational results of the genetic algorithm are discussed. The results are also compared with particle swarm optimization (PSO) in chapter 4. Conclusions of the paper and introduces possible future work are drawn in section 5.

2. Methods

2,1 Notation and Description of the PVRP

In general, periodic vehicle routing problems can be defined as a G = (C, E), where $C = (c_0, c_1, c_2, \dots, c_n)$ is the set that consists of subset $C = (c_1, c_2, \dots, c_n)$ which corresponds to the customers, and c_0 for the depot; and $E = \{d_{i,j} | c_{i,j} \in C, i \neq j\}$ is a set of distances, where $d_{i,j}$ represents the distance between c_i and c_j . A set of working days can be indexed by $D = (1, 2, \dots, T)$, where T denotes the total number of working days in a period length. The service demand of customer *i* can be denoted as r_i . Then let set $V = (v_1, v_2, \dots, v_k)$ as the vehicles at the depot, where *k* represents the total vehicles available at the depot. Additionally, the vehicle load capacity *v* is denoted as Q_v . As given by Chen, et al. [31], the mathematical model of the PVRP is as follows:

$$Objective = \min \sum_{i=0}^{n} \sum_{j=0}^{n} \sum_{t=1}^{T} \sum_{\nu=1}^{k} d_{i,j} z_{i,j,t,\nu}$$
(1)

Subject to:

$$\sum_{i=0}^{\infty} \sum_{j=0}^{\infty} z_{i,j,t,v} r_i \le Q_v \qquad \qquad i \ne j, v \in V, t \in D$$

$$\tag{2}$$

$$\sum_{\nu=1}^{n} z_{i,j,t,\nu} \le 1, i, j(i \ne j) \in N, t \in D$$

$$\tag{3}$$

 $\sum y_{t,v} \le k, t \in D \tag{4}$

$$\sum_{i,i,t,v}^{v=1} \in \{0,1\}, i, j \ (i \neq j) \in N, t \in D, v \in V$$
(5)

$$y_{t,v} \in \{0,1\}, t \in D, v \in V$$
 (6)

In this thesis, the objective of the PVRP is to minimize the total traveling time to all customers over the planning period, as expressed in Equation (1). Equation (2), ensures the total delivery size by vehicle v must not exceed the vehicle capacities on day t. Equation (3), make sure that the frequency of visits of vehicle v from customer i to customer j is limited to once a day at maximum. As shown in equation (4), ensure that the total number of vehicles assigned on day t $(y_{t,v})$ is restricted to the total vehicles available at the depot. Additionally, equation (5), $z_{i,j,t,v} = 1$ indicates vehicle v performs the services from

customer *i* to *j* on day *t*, while $z_{i,j,t,v} = 0$ indicating vice versa. Finally, in equation (6), $y_{t,v} = 1$ indicates that vehicle *v* is assigned to visit the customers on day *t*, while $y_{t,v} = 0$ indicating vice versa.

2.2 A Novel Coding Procedure

In this section, we provide a novel coding scheme to demonstrate the performance of each gene in the chromosome. The genetic algorithm has genetic operators, including crossover and mutation rates, which operate on the chromosome [32]. For a given chromosome that obtains by the encoding scheme, it is also required to know which permutation will correspond to [33]. This form of encoding allows genes that would help them to be constructed. Thus, the gene or chromosome is represented in this encoding scheme by a string of numbers representing its position in a sequence. The coding is based on the permutation of $\{1, 2, ..., n^* \max | b_i |\}$, where *n* describes the number of customers and max $| b_i |$ describes the total maximal replenishment times for customers (maximum trips). The procedure of the novel coding is as follows:

The types of customer visit frequency for six days are shown in Table 1. The sequence used represents a visit frequency of six days from Monday to Saturday, with 1 indicating a visit and 0 otherwise.

Visit frequency	Visit day combination		
visit irequency	1	2	3
2	100100	010010	001001
3	101010	010101	-
6	111111	-	-

Table	1. The Types of	of Customer	Visit Frequency

Based on customers' visit frequency, generate a random integer sequence from $\{1, 2, ..., n^* \max |b_i|\}$ to construct a feasible solution for PHVRP.

Compute $\{(Rand sequence + Normal sequence) \mod T\}+1$ to find the type of visit frequency at each customer, where T is the number of possible types for customers' visit frequency.

Based upon step (3), compute the combination of vehicles for each day with {(Total of the random sequence at each day $+ \mod K$)}+1, where K is the total type of vehicle combinations. The results imply as follow:

- One big vehicle is adopted if the result = 1.
- One small vehicle is adopted if the result = 2.
- One big and one small vehicle are adopted if the result = 3.
- Two big vehicles are adopted if the result = 4.
- Two small vehicles are adopted if the result = 5.

2.3 A Genetic Algorithm

Based on the novel coding scheme, we embedded the novel coding into a genetic algorithm. The GA can be described as a stochastic search algorithm inspired by natural competition between individuals to appropriate limited sources [34, 35]. The genetic algorithm effectively solves high computational complexity problems, including VRP. According to Rao, et al. [36], the computational procedure of the proposed GA to solve the PHVRP is classified as follows.

Define an appropriate chromosome to represent design parameters. Define the population in which consists of several groups of chromosomes, and the chromosomes are



Evaluate the fitness value or objective function. The fitness value in the genetic algorithm defines how close the obtained solution is to the optimal solution for the desired problem.

Apply the crossover rate and the mutation rate. The genetic algorithm needs to carry out the crossover and mutation rate operations to maintain the group's diversity and reduce errors. In this paper, we used Taguchi to find the best crossover and mutation rates, as shown in Table 2.

Algorithm
Initialization
Select parameters (crossover rate, mutation rate, affinity)
Generate an initial population based on the string number of the novel encoding
scheme
Repeat
while Total_iter<=Total_iter_max do
Evaluate the fitness of each individual in the population
for each individual
Calculate the fitness function
end
for i=1:length (Fitness), then sort generation by using
x = min (Fitness) - Fitness (i);
end
Then set as G _{bes}
Calculate the weight of each individual
$FIT = \frac{(Fitness - min(Fitness))}{((max)(Fitness))} + 0.0000001;$
((max(Fitness))-min(Fitness)) Maintenance of the fittest call before maturation retains good value
Croate a mating pool using a Boulotte wheel
Development the average operation
Powform the mutation of individuals
Further the new individuals
Evaluate the new individuals
End while
End while Until the store is so dition is so tisfied
Until the stopping condition is satisfied
Print the current best solution
Fig. 1. Procedure of genetic algorithm

The selection procedure was used to select genes from previously formed chromosomes. The Roulette Wheel Selection (RWS) method was one of the most commonly used selection methods. The parents obtained the chromosome's fitness value proportion using this method. The chromosome with the highest fitness value has been selected to have a higher probability than the other. Chromosome parents selected in the following selection process started the crossover process to produce offspring chromosomes. During the crossover process, the genotypes of both parents are merged to produce a new child.

Moreover, due to the high probability that the child will be exposed to further mutations with certain gene modifications. This aids in furthering the exploration of the solution area and ensuring or preserving genetic diversity. Meanwhile, the population must have the right balance of quality and genetic diversity to support efficient search.

Repeat step 3 until the maximum number of generations is achieved. The procedure of the genetic algorithm is presented in Fig. 1.

Table 2. Visit Frequency and Demand Quantity of the Customers						
Once per d	ay	Once per	2 days	Once per	3 days	
Node	Demand	Node	Demand	Node	Demand	
5	15	4	7	1	20	
13	10	8	10	2	22	
18	25	9	10	3	15	
22	27	16	15	6	12	
24	20	19	25	7	15	
		21	10	10	17	
		23	15	11	20	
				12	15	
				14	18	
				15	12	
				17	20	
				20	10	

2.4 Distribution Problem

The distribution problem considers the real-life instances of a distributor of liquefied petroleum gas (LPG) based in Batu, Indonesia. The distributor has 24 customers to be replenished frequently based on their given needs. Large and small vehicles with different capacities are available at the depot to serve every customer in each period. The working days of vehicles delivering the products are 6 days, from Monday to Saturday. The details of visit frequency and demand quantity of each customer are shown in Table 2. Some assumptions in this paper are as follows:

Each vehicle starts and terminates at the same depot. Each vehicle only visits each customer once after the scheduled delivery. Since the vehicle has a quantity, it can make multiple trips if the total quantity demand of all customers on day t exceeds its capacity.

There are five combinations of vehicles to be assigned each day: one big vehicle (M), one small vehicle (m), two big vehicles $(M_1 + M_2)$, two small vehicles $(m_1 + m_2)$, and one big vehicle and one small vehicle (M + m). The vehicles' capacity is set as 180 for the big vehicle and 60 for the small vehicle.

2.5 Experimental Setup

The genetic algorithm's initialization process needs to set some basic parameters. The correct choice of parameters will greatly affect the effectiveness of an algorithm. Setting different parameters of the Genetic Algorithm will produce different optimal results. Even if the same parameters are used in the algorithm, the results will still be different due to the randomness stochastic solution. The stochastic solution-search design is typical of the GA. The Taguchi method's main goal is to find optimal parameters to reduce this randomness. This paper provides the Taguchi Method to obtain the best parameter. Taguchi has been successfully applied in various experiments [37, 38]. The Taguchi model is based on optimizing key metrics known as signal-to-noise ratios (SNR) by conducting experiments using orthogonal arrays [39]. The mean square deviation of the objective function is called the signal-to-noise ratio (η). To calculate the stochastic

disturbance characteristics, each group of orthogonal table parameters was run 10 times for each experiment to represent the equivalence characteristics. Taguchi defined as,

$$\eta = -10 \log\left(\frac{1}{n} \sum_{t=1}^{n} y_t^2\right) \tag{7}$$

Equation (7) is Taguchi's as defined for *the smaller* the better characteristic. Where $\{y_1, y_2, ..., y_n\}$ denotes the objective function of experiments evaluation value. Thus, this paper adopted a three-level orthogonal array, a standard L_g (3²) array. The parameter combinations adopted in this paper are as follows: crossover rate (P_c) and mutation rate (P_m) as demonstrated by Hoogeboom, et al. [40], Mansour [41], and Ahmadi, et al. [42]. The overall computational result for the experiments is defined in Table 3. Thus, the result shows that the best parameter of crossover rate is 0.99 and 0.03 for mutation rate, with the lowest $\eta = -59.864$ indicating the optimal level.

The experimental results were performed using MATLAB R2015b software on Acer with specification Intel [®] Core [™] i7 – 3770 CPU [@] 3.40G Hz 64–bit and RAM 3.68 GB.

Fa	ctor	Objective						SNR				
P_{c}	P_m	y_1	y_2	y_3	y_4	y_5	<i>У6</i>	Y7	y_8	<i>у</i> 9	Y10	(η)
0.7	0.02	941.7	1030.7	946.5	936.8	935.6	1017.9	1063.7	1033.7	943.2	1006.2	-59.884
0.7	0.006	1039.1	1051.2	1031.1	942.8	1030.4	1037.6	987.0	943.6	934.4	1056.0	-60.055
0.7	0.03	1031.0	1035.4	1025.8	1017.6	931.5	1043.0	1041.4	1029.7	936.2	1109.0	-60.183
0.02	0.02	951.1	1030.1	1029.3	1030.2	1038.9	1028.3	1039.0	934.7	1021.2	1130.9	-60.211
0.02	0.006	1006.8	1019.5	1029.7	1033.2	1020.5	1032.2	1045.0	1048.3	1037.2	1039.1	-60.267
0.02	0.03	1105.7	1036.8	1079.5	1110.2	1045.0	1119.8	933.6	1122.8	937.5	1023.7	-60.453
0.99	0.02	1092.9	1044.1	1029.5	1032.9	1106.5	1035.4	936.5	1038.6	944.7	934.3	-60.182
0.99	0.006	938.3	1114.2	939.2	1050.8	938.7	1027.3	1052.5	1052.9	938.6	939.6	-60.011
0.99	0.03	1020.0	1018.9	1039.9	937.5	1034.3	937.0	933.0	1041.6	944.7	926.4	-59.864
Popul	Population = 100, generation = 2000, affinity = 0.9.											

Table 3. Numerical Result of Parameters of GA

3. Results and Discussion

This paper aims to minimize all customers' travel time over the planning period. Furthermore, the data related to the demand of each customer, the distant matrix, and the service schedule of each customer are collected in this research. Table 3 and Table 4 shows the computational and numerical results.

The solution of GA and PSO can be seen in Table 4. The combination of vehicles assigned to serve customers on day one is one big vehicle (M) with a total demand of 169 and a total traveling time of 141.5044 minutes for the GA. Similarly, the combination of vehicles assigned to serve the customer for day one of PSO is one big vehicle (M) with a total demand of 174 and a total traveling time of 147.4264 minutes. On day two, one big vehicle is assigned with a demand of 152 and a total traveling time of 144.3194 minutes for GA, while the total demand and traveling time of PSO are 150 and 163.0126, respectively. However, a combination of big and small vehicles (M+m) are assigned to serve customers on day three, with a loaded vehicle, 171 for big vehicles and 57 for small vehicles for GA. As well as PSO, the load of big vehicles is 173 and small vehicles is 57. The combination of vehicles to serve on day 4 is a big vehicle and a small vehicle (M+m) with the load vehicles for the big and small vehicles 124 and 47 for GA, and one big vehicle

with a loaded vehicle 169 for PSO. Next, on day five, one big vehicle (M) is assigned to serve the customer with a load is 155 for GA and 152 for PSO. Henceforward, the combination of big vehicle and small vehicle (M+m) are assigned to serve customers on day three with the load vehicles 171 for big vehicle and 59 for the small vehicle for GA. As well as PSO, the load of big vehicles is 173 and small vehicles is 57. Thus, the total traveling time is 1027.5960 minutes for GA and 1072.2484 minutes for PSO.

			GA		PSO		
Visit Day	Vehicle	Total Demand	Travelling Time (minutes)	Total Demand	Travelling Time (minutes)		
1	Big	169	141.5044	174	147.4264		
1	_	0	_	—	—		
9	Big	152	144.3194	150	163.0126		
2	—	0	—	—	—		
0	Big	171	164.4546	173	154.5005		
9	Small	57	63.0376	57	63.0376		
4	Big	124	142.2784	169	167.0226		
4	Small	47	15.4052	—	—		
F	Big	155	149.1824	152	147.8264		
9	_	—	—	_	—		
C	Big	171	149.4104	173	169.6786		
0	Small	59	58.0035	57	59.7436		
Total			1027.5960		1072.2484		

		\mathbf{GA}		PSO			
Description	Objective	CPU Time	Iteration	Objective	CPU Time	Iteration	
Min	997.23	578.28	364	1020.92	123.91	150	
Average	1001.52	590.65	1810	1079.25	125.09	764	
Max	1046.66	607.44	3685	1094.16	126.45	1582	
Standard Deviation	34.82	8.76	244.7	37.43	0.76	313.7	

Based on the result, it can be seen that the combination of the vehicle is only one big vehicle, and one big vehicle and one small vehicle. It implies that the company only requires one big vehicle and one small vehicle that should be available to deliver the product each day.

As shown in Table 5, the minimal objective of 100 experiments of GA is 997.23, and PSO is 1020.92, which means the min objective value of GA is superior to PSO. The computation time of the GA is longer than those of PSO, where the maximum CPU time of GA is 607.44, and PSO is 126.45. The PSO computation time is faster than GA because PSO can be implemented without many operators and parameters, such as crossover and mutation rates. GA can be implemented with many operators and parameters. GA also has higher computation efforts than PSO because PSO updates particles by themselves with the internal velocity and has a memory, which is important to the algorithm. The GA converges with minimal iterations is 364 and 150 for the PSO. The result implies that GA is outperformed, but PSO is superior in terms of iteration, computation time, and simplicity in finding the solution.



4. Comparison of Genetic Algorithm and Particle Swarm Optimization

To analyze the genetic algorithm, we also compared the objective result of the genetic algorithm and particle swarm optimization based on the average objective of numerical results with 30 experiments. We used the independent sample t-test to test further the genetic algorithm's performance and particle swarm optimization. We further test the following hypothesis with $\alpha = 0.025$.

The average objective of numerical results with 30 experiments can be seen in Table 6. As Shown in Table 7, the result of the t-test is -7.251 with a probability of 0.000, where less than α (0.025). Thus, it implies that the GA is significantly superior to PSO. Additionally, the average value of the GA is lower than that of the PSO. It can be concluded that the total traveling time of the GA is shorter than the total traveling time of the PSO, so it implies that GA outperforms the PSO.

	999.26	1031.37	1037.38	1019.28	1038.44	1021.22	1010.38	1024.67	1039.20	1044.97
GA	1010.77	1036.62	999.01	1040.21	1040.47	1019.44	1037.94	1022.35	1011.15	1001.96
	1004.45	1044.28	1022.87	1041.31	1007.88	1025.68	1007.38	1005.05	1019.51	1036.32
	1024.49	1074.16	1059.80	1051.28	1044.01	1084.68	1057.59	1073.73	1059.28	1033.46
PSO	1042.24	1050.52	1079.64	1032.02	1081.28	1058.32	1035.16	1034.48	1051.94	1074.88
	1093.99	1078.25	1021.91	1048.17	1058.98	1035.62	1047.63	1092.13	1046.46	1086.05

Table 6. The average objective of numerical results with 30 experiments

Table 7	The resul	t of the	statistical	test of	GA and PSO
Table 1.	THE LESU	U UI UIIC	statistical		UL and 1 DO

	Methods	Mean	t	P Value
GA		1023.36	7 051	0.000
PSO		1057.07	7.201	0.000

5. Conclusion

This paper has proposed a genetic algorithm to solve the periodic heterogeneous vehicle routing problem (PHVRP). The objective is to minimize the total traveling time during the distribution. Using an orthogonal array, we used the Taguchi method to find the best crossover rate and mutation rate of the genetic algorithm. The real case of PHVRP in Batu City, Indonesia has been solved using the genetic algorithm. Furthermore, the numerical results show that PSO's objective value is greater than GA's. However, the computation time is faster than those of GA. Additionally, the statistical test show that the average objective of GA is superior to those of PSO. In future research, a possible algorithm improvement could also involve comparing results with the proposed algorithm.

Declarations

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